## Claims

We claim:

1	1. A system for organizing and accessing a database, the system comprising:
2	a primary B+tree index;
3	a secondary B+tree index;
4	a plurality of mapping table row identifiers stored in the secondary B+tree index; and
5	a plurality of database addresses for leaf blocks of the primary B+tree index
6	corresponding to mapping table rows.
1	2. The system according to claim 1, wherein 4 bytes of the database addresses are stored
2	in the secondary B+tree index.
1	3. The system according to claim 1, wherein the database addresses correspond to a
2	guess-database address.
1	4. The system according to claim 1, further comprising:
2	a guess-database address quality statistic for the secondary index and a guess-database
3	address quality statistic for the mapping table, both statistics being operable to assess guess-
4	database address quality.
1	5. A method for managing a database system, the method comprising:

2	creating a secondary index for a B+tree structure, wherein the secondary index structure
3	comprises a plurality of rows each comprising an index key value, a mapping table rowid value
4	and a guess-database address value.
1	6. The method according to claim 5, further comprising:
2	inserting a row of the secondary index structure, wherein inserting the row comprises
3	inserting a row comprising an index key value, a mapping table rowid value and a guess-
4	database address value.
1	7. The method according to claim 5, further comprising:
2	deleting a row of the secondary index, wherein deleting the row comprises locating a row
3	comprising an index key value and a mapping table row identifier and deleting the row.
1	8. The method according to claim 5, further comprising:
2	updating the secondary index, wherein updating the secondary index comprises locating a
3	row of the secondary index comprising an old index key and a mapping table row identifier,
4	deleting the row and inserting in the row a new index key value, a mapping table row identifier
5	and a guess database address.
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1	9. The method according to claim 5, further comprising carrying out a query utilizing the
2	secondary index, wherein carrying out the query comprises:
3	utilizing a guess-database address stored as part of a secondary index row to find a row in
4	the primary B+tree structure;

)	obtaining a target database block from the row in the primary B+tree structure;
6	searching the target database block for a row that contains a mapping table row identifier
7	that is the same as a mapping table row identifier stored in the secondary index row; and
8	if the row in the database block matches the target database block, then the correct row in
9	the database has been located and the query is completed.
1	10. The method according to claim 9, wherein if the row in the database block does not
2	match the target database block, carrying out the query further comprises:
3	accessing the mapping table row stored in the secondary index row;
4	utilizing a guess-database address stored in the mapping table row to access a target block
5	of the database;
6	searching the target block for a primary key that matches a primary key stored in the
7	mapping table row; and
8	if the primary key is found, then the query is completed.
1	11. The method according to claim 10, wherein if the primary key is not located carrying
2	out the query further comprises:
3	traversing the primary B+tree structure utilizing the primary key value from the mapping
4	table row to identify the database address to complete the query.
1	12. The method according to claim 11, further comprising:
2	maintaining a guess-database address quality statistic for the secondary index;
3	maintaining a guess-database address quality statistic for the mapping table;

4	utilizing the statistics to assess guess-database address quality; and
5	carrying out the query based upon guess-database quality in the secondary index and
6	mapping table.
1	13. The method according to claim 12, further comprising:
2	estimating guess-database address quality;
3	estimating the cost of the query based upon the estimated guess-database address quality;
4	and
5	carrying out the query starting with an index structure with the lowest estimated cost.
1	14. The method according to claim 5, wherein only 4 bytes of the guess-database address
2	value are stored in the secondary index row.
1	15. A computer program product for performing a process of managing a database system,
2	comprising:
3	a computer readable medium; and
4	computer program instructions, recorded on the computer readable medium, executable
5	by a processor, for performing the steps of:
6	creating a secondary index for a B+tree structure, wherein the secondary index structure
7	comprises a plurality of rows each comprising an index key value, a mapping table rowid value
8	and a guess-database address value.

16. A system for performing a database management process, comprising

a processor operable to execute computer program instructions; and
a memory operable to store computer program instructions executable by the processor,
for performing the steps of:
creating a secondary index for a B+tree structure, wherein the secondary index structure
comprises a plurality of rows each comprising an index key value, a mapping table rowid value

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and a guess-database address value.